New Data Storage Model for H1

Thomas Benisch, DESY

CHEP 2000 February 7–11, 2000 Padova, Italy

Introduction

Shortcomings of the present Data Storage Model

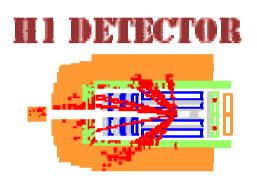
New Data Storage Model

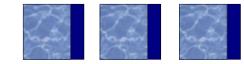
Benchmark Results

Conclusions

Introduction (1)

- H1 is an experiment at the ep collider HERA at DESY in Hamburg (Germany)
- □ H1 started collecting data in 1992
- H1 and HERA are performing substantial upgrades in the year 2000
 - increased demands on data storage and data handling





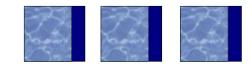
Introduction (2)

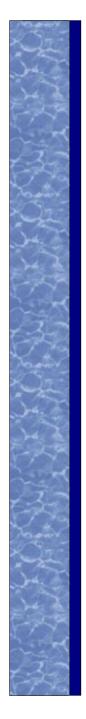
the H1 collaboration has decided to adopt the ROOT framework for

- physics analysis
- event display

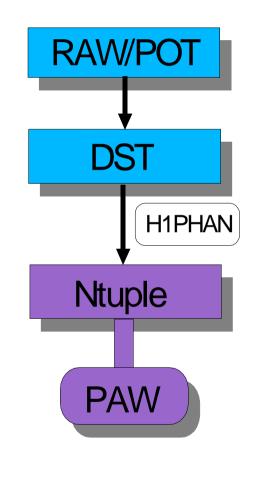
 \Rightarrow this talk

 \Rightarrow see talk of U. Berthon





Present Data Storage Model

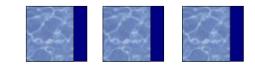


- RAW Raw Data
- POT Production Output Tape raw + reconstructed events
- DST Data Summary Tape reduced set of information for physics analysis

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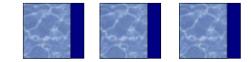
BOS/FPACK

ZEBRA



Data Structure and Data Access

- data storage is based on
 - memory management system BOS (Bank Organisation System)
 - □ I/O package FPACK
- Entity–Relationship model is used as a basis for data structures
- data are stored in so-called BOS banks
 - described by a data definition language (DDL)
 - contain closely related data
- I/O package FPACK
 - machine independent
 - common user interface for all operating systems
 - network server for remote file access



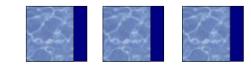
Shortcomings of the Present System

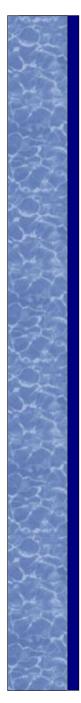
- The current H1 analysis environment is rather inhomogenous
 - data storage based on BOS/FPACK
 - physics analysis is done with PAW
 - event display uses graphics package LOOK
- events must always be read and written completely, reading of single variables is not possible
- event selection is done by using index files, based on a 32 bit classification word (predefined, static!)



New Technology for Data Storage: ROOT

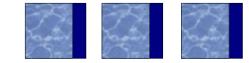
- increased demands on data handling after the upgrade in the year 2000
- H1 has chosen a new technology to improve data access: ROOT
- important features of ROOT I/O:
 - high I/O efficiency
 - sequential and direct data access techniques
 - possibility of splitting event data into several streams
 - access to part of the event data only
 - support of networking
 - built-in gzip-type compression algorithm

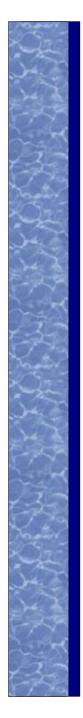




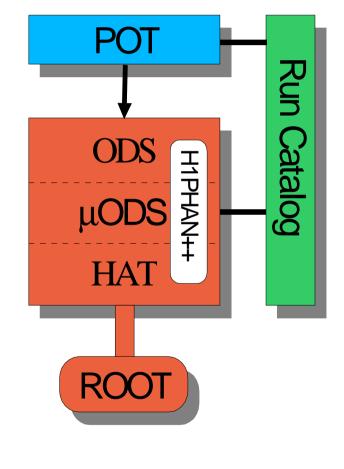
New Data Storage Model (1)

- event data are split into several parts according to frequency of their access
- multi-level hierarchical storage system
- higher level contains less data per event
- access to partial event data will be essential for many physics analyses





New Data Storage Model (2)

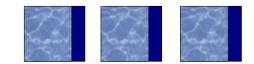


POT	Production	Output	Tape
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- ODS Object Data Store
- **μODS** micro Object Data Store
- □ HAT H1 Analysis Tag

BOS/FPACK

ROOT



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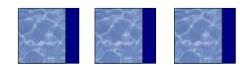
New Data Storage Model (3)

POT Production Output Tape **BOS/FPACK**

- contains raw and reconstructed data as now
- no change to software producing the POTs necessary, e.g. Fortran based reconstruction program

ODS Object Data Store ROOT

- contains all standard objects for physics analysis, e.g. tracks, clusters, etc.
- corresponds to current DSTs
- existing event model is not changed
- direct mapping of BOS banks into objects using BOS200P
- each bank is written to a seperate branch
- backward compatibility



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PO

ODS HIPHANH

ataloc

ODS

HAT

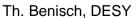
ROO

New Data Storage Model (4)

micro Object Data Store μODS ROOT

- contains information on the particle level and their 4-vectors, e.g. electron candidates, jet properties, energy flow, PID probabilities, etc.
 - full expert knowledge of all physics working groups
 - filling code will be part of new physics analysis package H1PHAN++
 - will be written in C++ in an object-oriented way
 - based on algorithms of old physics analysis package





PO

ODS

HAT

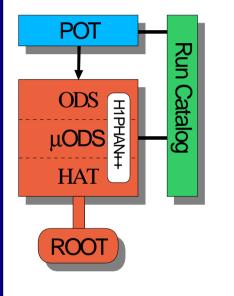
ROO

ODS HIPHANH

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New Data Storage Model (5)

HAT H1 Analysis Tag ROOT

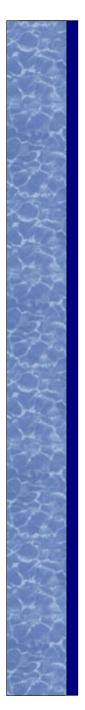


- □ tag database for fast and dynamic event selection
- contains about 100 kinematic event variables, e.g. Q², x, y, etc.
- experience from present tag database (Objectivity/DB)
- Run Catalog

- MySQL/Oracle
- access to different parts of the event
- prototype will be implemented in MySQL
- Oracle as alternative



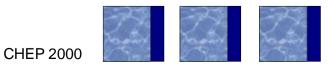
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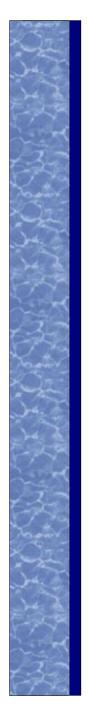


Data Volumes

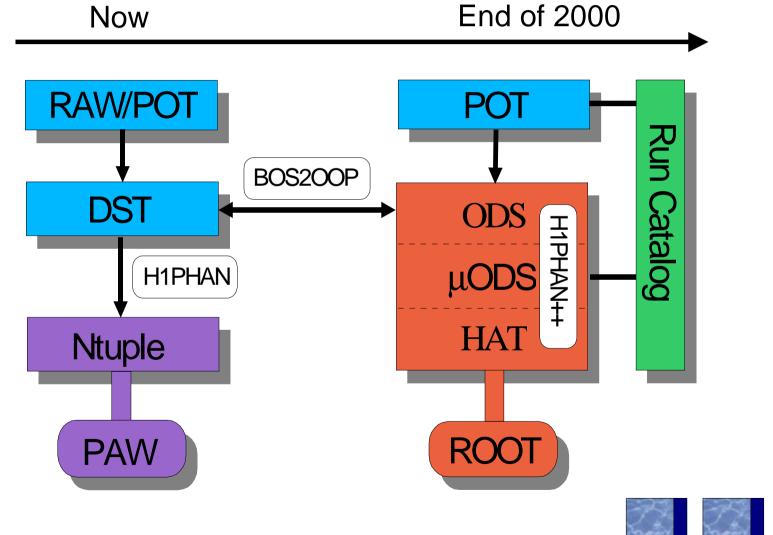
H1 data volumes after the luminosity upgrade in the year 2000

storage level	storage media	event size (kB)	data volume (GB/year)
РОТ	tape	200	10000
ODS	disk	15	500
μ ODS	disk	3	100
НАТ	disk	0.5	15

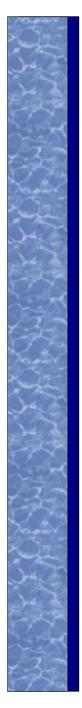




Timeline for Implementation



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Benchmark Results (1)

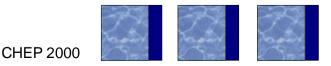
□ ROOT I/O performance comparisons were already done by

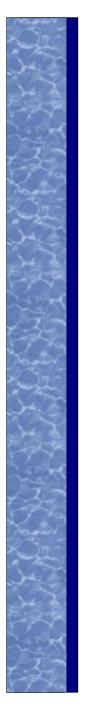
- ROOT team
 - ROOT versus ZEBRA
 - ROOT versus Objectivity/DB
- CDF
 - ROOT versus YBOS

 \Rightarrow significant speed-up in data access time with ROOT I/O

Benchmark Results (2) H1 compared ROOT and FPACK convert FPACK DST file into ROOT format ROOT file factor 3 smaller for compression level 2 ROOT file factor 1.3 bigger for no compression

- read whole event data
 - ROOT about a factor 4.6 slower than FPACK
 - no optimization done yet
- read only part of the event data
 - no results from H1 available yet, but significant speed improvements expected





Conclusions

the new H1 data storage model has been presented

 the choice of ROOT will provide an efficient solution for data handling and storage after the upgrade in the year 2000

